

# ECE 2574: Data Structures and Algorithms - Applications of Recursion I

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Today we will look a common task that is easily solved using a recursive solution, parsing algebraic expressions.

- ▶ Warmup
- ▶ Review of algebraic expressions
- ▶ converting between prefix and postfix expressions
- ▶ an implementation in C++

## Why prefix or postfix notation?

Many of you asked, in some form or another, why prefix or postfix notation?

- ▶ unambiguous, there are no operator precedence rules
- ▶ easy to parse (translate into a tree form) for evaluation
- ▶ supports operators with n-ary arguments with no additional syntax

For this reason, prefix and postfix notation is used in many programming and data description languages

- ▶ languages in the Lisp family use prefix notation
- ▶ stack-based languages generally use postfix notation

# Algebraic Expressions

Lets say we are going to write a program to act as a calculator. For example:

$$(a + b)*c$$

$$(a/b)*c$$

$$(a-b-c-d)/e$$

How does the calculator decide if the expression is valid ?

## Lets start with a less complex Algebraic grammar

**Prefix expressions.** In prefix notation the operation is written first, followed by the two operands.

Examples:

▶ \* + a b c in infix notation is  $(a + b) * c$

▶ + / a b - c d in infix notation is  $(a / b) + (c - d)$

The grammar looks like:

$\langle \text{prefix} \rangle = \langle \text{operand} \rangle \mid \langle \text{operator} \rangle \langle \text{prefix} \rangle \langle \text{prefix} \rangle$

$\langle \text{operator} \rangle = + \mid - \mid * \mid /$

$\langle \text{operand} \rangle = a \mid b \mid c \mid \dots \mid z$

## Validation of a prefix expression using recursion

$\langle \text{prefix} \rangle = \langle \text{operand} \rangle \mid \langle \text{operator} \rangle \langle \text{prefix} \rangle \langle \text{prefix} \rangle$   
 $\langle \text{operator} \rangle = + \mid - \mid * \mid /$   
 $\langle \text{operand} \rangle = a \mid b \mid c \mid \dots \mid z$

Base step is simple: check for operator at string beginning.

The recursive step is a little more complicated. The key is that if  $\langle \text{prefix} \rangle$  is a valid prefix  $\langle \text{prefix} \rangle \langle \text{ch} \rangle$ , where  $\langle \text{ch} \rangle$  is any non-blank character, is not.

## Validation of a prefix expression using recursion

```
function endPre(in s:string, in first:int): int

last = s.length() - 1
if( first < 0 or first > last )
    return -1

ch = first char of s
if(ch is an operand)
    return first
else if(ch is an operator)
    firstEnd = endPre(s, first+1);
    if(firstEnd > -1)
        return endPre(s, firstEnd +1)
    else
        return -1
else
    return -1
endfunction
```

## Using the endPre function to validate the grammar

Call endPre at first character in the string  
if the last character returned is not the last one it is not a valid  
prefix expression.

```
function isPre(in s:string): bool
```

```
  lastChar = endPre(s, 0)
```

```
  return  lastChar >= 0 AND  
         lastChar == s.length()-1
```

```
endfunction
```



## Warmup #1

Is the following string a valid prefix expression?

/ + a c d - e g

False (53% correct)

## Similar is the postfix notation

`<postfix> = <operand> | <postfix> <postfix>`  
`<operator> <operator> = + | - | * | /`  
`<operand> = a | b | c | .... | z`

Suppose we wanted to convert the prefix expression to a postfix expression.

`<postfix> = <operand> | <postfix> <postfix> <operator>`  
`<prefix> = <operator> | <operator> <prefix> <prefix>`

## A recursive solution to conversion

```
function convert(in pre:string,
                out post:string)
  ch = first character of pre
  delete first character of pre

  if ch is an operand
    post = post + ch //concatenate
  else
    // recursion to convert 1st
    convert(pre, post)
    // recursion to convert 2nd
    convert(pre, post)
    // concatenate the operator
    post = post + ch
  endif

endfunction
```

## Validating a postfix expression

```
function endPost(in s:string, in last:int): int

first = 0
if( first > last )
    return -1

ch = last char of s
if(ch is an operand)
    return last
else if(ch is an operator)
    lastEnd = endPost(s, last-1);
    if(lastEnd > -1)
        return endPost(s, lastEnd-1)
    else
        return -1
else
    return -1
endfunction
```

## Validating a postfix expression: isPost

Call endPost at last character in the string  
if the last position returned is not zero it is not a valid postfix  
expression.

```
function isPost(in s:string): bool  
  
firstChar = endPost(s, s.length()-1)  
  
return firstChar == 0  
  
endfunction
```

## Warmup #2

Is the following string a valid postfix expression?

h r \* R f - + t g - e f / \*

False (49% Correct)

## A recursive solution to conversion the other way

```
function convert(in post:string,
                out pre:string)
  ch = last character of post
  delete last character of post

  if ch is an operand
    pre = pre + ch //concatenate
  else
    // concatenate the operator
    pre = pre + ch
    // recursion to convert 1st to temp
    convert(post, temp)
    // recursion to convert 2nd
    convert(post, pre)
    pre = pre + temp // append temp
  endif
endfunction
```

## Warmup #3

Convert the following prefix expression to a postfix expression.

+ \* A B / C D

A B \* C D / + (80 Correct)



## Exercise: Implementing and testing in C++

See website.

## Next Actions and Reminders

- ▶ Read CH pp. 172-186.
- ▶ There is no warmup for Monday.
- ▶ Program 1 is due Wed at 11:55pm **via Canvas**.
- ▶ If you have used all your late days, you **must** turn it in on time.