

ECE 2574: Data Structures and Algorithms - Recursion Part I

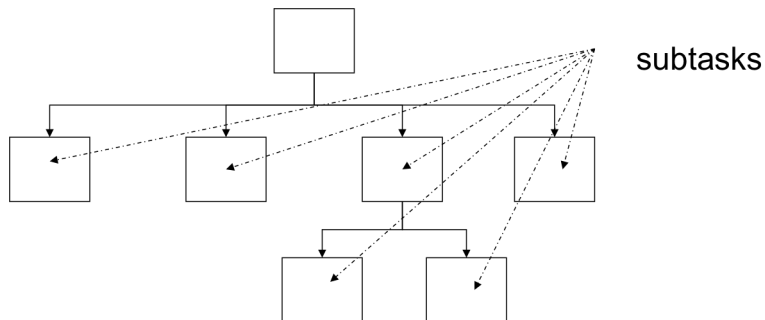
C. L. Wyatt

Today we will introduce the notion of recursion, look at some examples, and see how to implement them in code.

- ▶ Introduction to recursion
- ▶ Warmup
- ▶ Examples
- ▶ Exercise: Recursive Egyptian Powers

Top-down Design

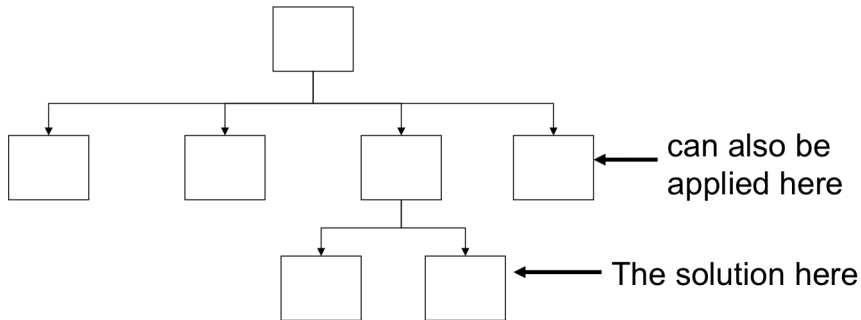
Top down design divides problems into smaller and easier sub-problems.



The hope is, at each of these successive levels, these sub-problems are easier to solve.

Recursion

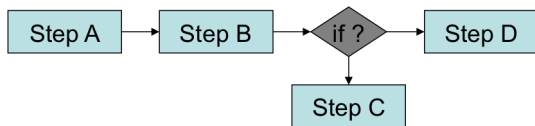
In some cases, the solution to lowest level sub-problem can be applied at the higher level.



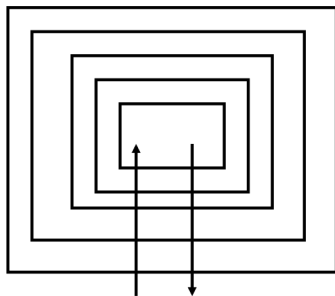
This type of solution is called *recursive*.

A graphical view

Graphical depictions of algorithms show connected boxes.



Recursive solutions form nested boxes.



- ▶ The inner-most box solves the lowest-level problem,
- ▶ The next box solves the next level above.
- ▶ In a recursive solution the algorithm is the same at each level.

Example: $n!$, the factorial

Simple iterative solution to compute the factorial

```
int factorial (int n)
{
    int result = n;
    do
    {
        n -= 1;
        result = result*n;
    } while (n > 1);

    return result;
};
```

How large an n will this work for in a real programming languages?

We can break the factorial solution into a recursive solution.

$$n! = 1*2*3*4*\dots*n$$

grouping terms

$$n! = 1*2*3*4*\dots*(n-1)*n$$

$$n! = ((n-1)!) * n$$

This is an example of a recurrence relation.

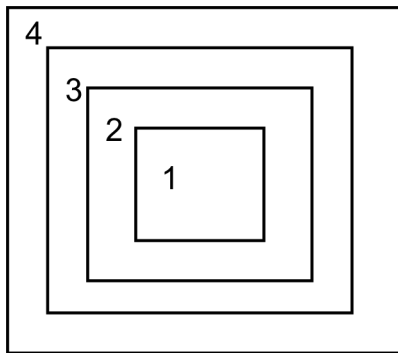
A recursive solution to the factorial

```
int factorial (int n)
{
    if (n <= 1)
        return 1;
    else
        return(n*factorial(n-1));
};
```

A recursive solution to the factorial

Graphically each box takes the output of the box inside it and multiplies by the integer in that box.

Example: $4! = 4 * 3 * 2 * 1 = (4 * (3 * (2 * (1))))$



Formal definition of recursion

- ▶ A recursive procedure is one whose evaluation at (non-initial) inputs involves invoking the procedure itself at another input.
- ▶ In the case of the factorial this involves invoking the procedure at (n-1) when the input is n: $n! = n \cdot (n-1)!$
- ▶ Recursion is a very powerful tool in the design and analysis of algorithms.
- ▶ Often complex problems have very simple recursive solutions.

What makes recursive procedures work ?

* At each invocation, the solution must get closer to a known solution

i.e. $0! = 1! = 1$

- ▶ The procedure calls must terminate in a finite number, that is the function must not endlessly call itself. Otherwise the recursion is *infinite*

Recursive version of the GCD algorithm

Recall the GCD algorithm

- ▶ A.0 If $m < n$, swap m and n .
- ▶ A.1 Divide m by n and let r be the remainder.
- ▶ A.2 If $r = 0$, terminate; n is the answer.
- ▶ A.3 Set m to n , n to r , and go back to step A.1.

A recursive solution (after step 0).

```
int gcd(int m, int n)
{
    if( n == 0 ) return m;
    else return gcd(n, m%n);
}
```

Recursive version of the GCD algorithm

The recurrence relation for GCD

$$\text{gcd}(m,n) = \text{gcd}(n, m \bmod n)$$

The stopping condition (base case) is $n = 0$

Example:

$$\begin{aligned} \text{gcd}(131,62) &\rightarrow \text{gcd}(62,7) \rightarrow \text{gcd}(7,6) \\ &\rightarrow \text{gcd}(6,1) \rightarrow \text{gcd}(1,0) \end{aligned}$$

Warmup #1

Which of the following C++ functions correctly computes the sum from 1 to n using recursion?

```
int function1(int n){
    int sum = 0;
    for(int i = 1; i <= n; ++i){
        sum += i;
    }
    return sum;
}
```

Incorrect. (15%).

Warmup #1

Which of the following C++ functions correctly computes the sum from 1 to n using recursion?

```
int function2(int n){  
    if(n == 1) return 1;  
    return n + function2(n-1);  
}
```

Correct (81%).

Warmup #1

Which of the following C++ functions correctly computes the sum from 1 to n using recursion?

```
int function3(int n){  
    return n*(n+1)/2;  
}
```

Incorrect. (3%).

Warmup #2

What would happen if function1 in the previous question was called with an argument of -1?

```
int function1(int n){
    int sum = 0;
    for(int i = 1; i <= n; ++i){
        sum += i;
    }
    return sum;
}
```

The correct answer is “it would return 0” (63%).

Warmup #3

What would happen if function2 in the previous question was called with an argument of -1?

```
int function2(int n){  
    if(n == 1) return 1;  
    return n + function2(n-1);  
}
```

The correct answer is “A run-time error would occur” (73%).

Warmup #4

What would happen if function3 in the previous question was called with an argument of -1?

```
int function3(int n){  
    return n*(n+1)/2;  
}
```

The correct answer is “It would return 0” (85%).

Another Example: Exponentiation

- ▶ The Egyptian Powers algorithm computes x to the power n by repeated squaring.
- ▶ The recurrence relation for computing x^n for any positive integer n :

$$x^n = \begin{cases} (x \cdot x)^{n/2} & n \text{ even} \\ x(x \cdot x)^{(n-1)/2} & n \text{ odd} \end{cases}$$

Exercise: write a recursive function and a set of tests implementing the Egyptian Powers algorithm.

In pseudo-code

```
function RecPowers (x, n)
```

Input: a real number x and positive integer n

Output: x raised to power n

```
    if (n == 1) // initial condition
```

```
        pow = x;
```

```
    else
```

```
        if even(n) then
```

```
            pow = RecPowers(x*x,n/2)
```

```
        else
```

```
            pow = x*RecPowers(x*x,(n-1)/2)
```

```
        endif
```

```
    endif
```

```
    return (pow)
```

```
endfunction
```

Next Actions and Reminders

- ▶ Read CH pp. 67-87
- ▶ Warmup before noon on Monday.
- ▶ Program 0 due tonight by 11:59 PM.